

Canonical gradings of monads

Flavien Breuvert Dylan McDermott Tarmo Uustalu

Example

The writer monad Wr for lists over a set C has:

object mapping $\text{Wr} : \mathbf{Set} \rightarrow \mathbf{Set}$ $\text{Wr } X = \text{List } C \times X$

unit functions $\eta_X : X \rightarrow \text{Wr } X$ $\eta_X x = ([], x)$

multiplication functions $\mu_X : \text{Wr } (\text{Wr } X) \rightarrow \text{Wr } X$ $\mu_X (s_1, (s_2, x)) = (s_1 \# s_2, x)$

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$$\text{multiplication functions} \quad \mu_X : \text{Wr} (\text{Wr } X) \rightarrow \text{Wr } X \quad \mu_X (s_1, (s_2, x)) = (s_1 \# s_2, x)$$

We can grade this by

- ▶ natural numbers $e \in \mathbf{N}$, to get a graded monad WrL :

$$\text{WrL } e X = \text{List}_{\leq e} C \times X \quad \eta : X \rightarrow \text{WrL } 0 X \quad \mu : \text{WrL } e_1 (\text{WrL } e_2 X) \rightarrow \text{WrL } (e_1 + e_2) X$$

- ▶ subsets $e \subseteq C$, to get a graded monad WrS :

$$\text{WrS } e X = \text{List } e \times X \quad \eta : X \rightarrow \text{WrS } \emptyset X \quad \mu : \text{WrS } e_1 (\text{WrS } e_2 X) \rightarrow \text{WrS } (e_1 \cup e_2) X$$

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We can grade this by

- ▶ subsets $\Sigma \subseteq \text{List } C$, to get a graded monad WrC :

$$\text{WrC } \Sigma X = \Sigma \times X \quad \eta : X \rightarrow \text{WrC } J X \quad \mu : \text{WrC } \Sigma_1 (\text{WrC } \Sigma_2 X) \rightarrow \text{WrC} (\Sigma_1 \boxplus \Sigma_2) X$$

where

$$J = \{[]\} \quad \Sigma_1 \boxplus \Sigma_2 = \{s_1 \# s_2 \mid s_1 \in \Sigma_1, s_2 \in \Sigma_2\}$$

Example

WrC is the canonical grading of Wr:

- ▶ WrL is

$$\mathbf{N} \xrightarrow{F} \mathcal{P}(\text{List } C) \xrightarrow{\text{WrC}} [\mathbf{Set}, \mathbf{Set}]$$

where

$$Fe = \text{List}_{\leq e} C \subseteq \text{List } C$$

- ▶ WrS is

$$\mathcal{P}C \xrightarrow{F} \mathcal{P}(\text{List } C) \xrightarrow{\text{WrC}} [\mathbf{Set}, \mathbf{Set}]$$

where

$$Fe = \text{List } e \subseteq \text{List } C$$

This work

For every monad T on **Set**:

- ▶ there is a notion of grading of T
- ▶ T has a canonical grading
- ▶ every other grading factors through the canonical one

This work

More generally, given a suitable notion of \mathcal{M} -subfunctor, for every monad T :

- ▶ there is a notion of \mathcal{M} -grading of T
- ▶ T has a canonical \mathcal{M} -grading
- ▶ every other \mathcal{M} -grading factors through the canonical one

In particular, monads can be canonically graded by shapes

Gradings

Let T be a monad on \mathbf{Set}

$$T : \mathbf{Set} \rightarrow \mathbf{Set} \quad \eta : X \rightarrow TX \quad \mu : T(TX) \rightarrow TX$$

A **grading** (\mathcal{E}, G) of T consists of:

- ▶ an partially ordered monoid $(\mathcal{E}, \leq, I, \odot)$ of **grades** $e \in \mathcal{E}$
- ▶ a subset $GeX \subseteq TX$ for each $e \in \mathcal{E}$ and set X

such that

- ▶ $GeX \subseteq Ge'X$ for all $e \leq e'$ and X
- ▶ G is closed under the monad structure of T :

$$Tf : TX \rightarrow TY \text{ restricts to } Gef : GeX \rightarrow GeY \quad \text{for each } e \in \mathcal{E} \text{ and } f : X \rightarrow Y$$

$$\eta : X \rightarrow TX \text{ restricts to } \eta : X \rightarrow GIX \quad \text{for each } X$$

$$\mu : T(TX) \rightarrow TX \text{ restricts to } \mu : Ge_1(Ge_2X) \rightarrow G(e_1 \odot e_2)X \quad \text{for each } e_1, e_2 \in \mathcal{E} \text{ and } X$$

Fact: (G, η, μ) is a graded monad

Gradings

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such that some conditions hold.

- ▶ $(\mathcal{E}, \leq, I, \odot) = (\mathbf{N}, \leq, 0, +)$, and

$$\text{WrL } eX = \text{List}_{\leq e} C \times X \subseteq \text{List } C \times X = \text{Wr } X$$

- ▶ $(\mathcal{E}, \leq, I, \odot) = (\mathcal{P}C, \subseteq, \emptyset, \cup)$, and

$$\text{WrS } eX = \text{List } e \times X \subseteq \text{List } C \times X = \text{Wr } X$$

- ▶ $(\mathcal{E}, \leq, I, \odot) = (\mathcal{P}(\text{List } C), \subseteq, \emptyset, \cup)$, and

$$\text{WrS } eX = \text{List } e \times X \subseteq \text{List } C \times X = \text{Wr } X$$

Aside: constructing the ordered monoid

Given

- ▶ a set \mathcal{E}
- ▶ a subfunctor $Ge \subseteq T$ for each $e \in \mathcal{E}$

under some conditions we get a grading of T by defining

$$e \leq e' \iff Ge \subseteq Ge'$$

$$I = \text{smallest } e \text{ such that } \eta \text{ restricts to } \eta : X \rightarrow GeX$$

$$e_1 \odot e_2 = \text{smallest } e \text{ such that } \mu \text{ restricts to } \mu : Ge_1(Ge_2X) \rightarrow GeX$$

Canonical grading of a functor

A **grading** (\mathcal{E}, G) of a functor T consists of:

- ▶ a partially ordered set (\mathcal{E}, \leq) of **grades** $e \in \mathcal{E}$
- ▶ a subfunctor $Ge \subseteq T$ for each $e \in \mathcal{E}$

such that $Ge \subseteq Ge'$ for all $e \leq e'$

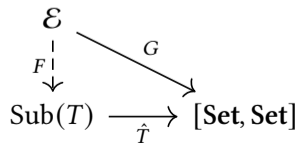
The **canonical** grading $(\text{Sub}(T), \hat{T})$ of T has:

(ignoring some size issues)

poset	$\text{Sub}(T)$	(subfunctors of T , ordered by pointwise inclusion)
subfunctors	$\hat{T}S = S$	

Universal property:

for every other grading (\mathcal{E}, G) of T , there is a unique monotone $F : \mathcal{E} \rightarrow \text{Sub}(T)$ such that $Ge = \hat{T}(Fe)$ for all e



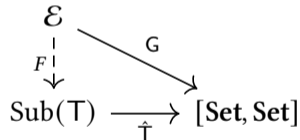
Canonical grading of a monad

The canonical grading $(\text{Sub}(T), \hat{T})$ of a monad T has

ordered monoid $(\text{Sub}(T), J, \boxplus)$ where $JX = \{\eta x \mid x \in X\}$
 subfunctors $\hat{T}S = S$ where $(S_1 \boxplus S_2)X = \{\mu t \mid t \text{ is in the image of } S_1(S_2X) \hookrightarrow T(TX)\}$

Universal property:

for every other grading (\mathcal{E}, G) of T , there is a unique $F : \mathcal{E} \rightarrow \text{Sub}(T)$ that is lax monoidal and satisfies $Ge = \hat{T}(Fe)$ for all e



$e \leq e' \Rightarrow Fe \subseteq Fe'$ $J \subseteq FI$ $Fe \boxplus Fe' \subseteq F(e \odot e')$

Example: writer

Take the writer monad Wr

$$\text{Wr } X = \text{List } C \times X$$

Subfunctors $S \subseteq \text{Wr}$ are equivalently subsets

$$\Sigma \subseteq \text{List } C$$

via

$$\Sigma = \{s \in \text{List } C \mid (s, \star) \in S1\}$$

$$SX = \{(s, x) \in \text{List } C \times X \mid s \in \Sigma\}$$

So the canonical grading is $\mathcal{P}(\text{List } C)$ with

$$\text{Wr } C \Sigma X = \Sigma \times X$$

$$\Sigma \leq \Sigma' \Leftrightarrow \Sigma \subseteq \Sigma' \quad J = \{[]\}$$

$$\Sigma_1 \boxplus \Sigma_2 = \{s_1 \# s_2 \mid s_1 \in \Sigma_1, s_2 \in \Sigma_2\}$$

$$\begin{array}{ccc} \mathbf{N} & & \\ \downarrow F & \searrow \text{WrL} & \\ \mathcal{P}(\text{List } C) & \xrightarrow{\text{WrC}} & [\text{Set}, \text{Set}] \end{array}$$
$$Fe = \{s \in \text{List } C \mid |s| \leq e\}$$

$$\begin{array}{ccc} \mathcal{P}C & & \\ \downarrow F & \searrow \text{WrS} & \\ \mathcal{P}(\text{List } C) & \xrightarrow{\text{WrC}} & [\text{Set}, \text{Set}] \end{array}$$
$$Fe = \text{List } e$$

Example: reader


Take the reader monad Read_V (for a set V)

$$\text{Read}_V X = V \Rightarrow X$$

Subfunctors $S \subseteq \text{Read}_V$ are equivalently upwards-closed sets

$$\Sigma \subseteq \text{Equiv}_V$$

$R \in \Sigma \Rightarrow R' \in \Sigma \text{ whenever } R \subseteq R'$



of equivalence relations of V , via

$$\Sigma = \{R \in \text{Equiv}_V \mid [-]_R \in S(V/R)\}$$

$$SX = \{f : V \rightarrow X \mid \exists R \in \Sigma. \forall v, v'. v R v' \Rightarrow f v = f v'\}$$

and these give a canonical grading $(\text{Sub}(\text{Read}_V), \text{ReadC}_V)$

Example: for $FR = \{R' \in \text{Equiv}_V \mid R \subseteq R'\}$, the graded monad

$$\text{Equiv}_V \xrightarrow{F} \text{Sub}(\text{Read}_V) \xrightarrow{\text{ReadC}_V} [\mathbf{Set}, \mathbf{Set}]$$

is

$$\text{Read}'_V R X \cong (V/R) \Rightarrow X$$

\mathcal{M} -gradings of functors

For a class \mathcal{M} of natural transformations \succrightarrow , an \mathcal{M} -grading of a functor T consists of

- ▶ a category \mathcal{E}
- ▶ a functor $G : \mathcal{E} \rightarrow [\text{Set}, \text{Set}]$
- ▶ a natural transformation $m_e : Ge \succrightarrow T$, whose components are in \mathcal{M}

The \mathcal{M} -subfunctors of T form an \mathcal{M} -grading $(\mathcal{M}/T, \hat{T})$, with

$$\hat{T}(S, m) = S \quad \succ \xrightarrow{m} \quad T$$

and this is canonical:

$$\begin{array}{ccc} \mathcal{E} & & \\ \downarrow F & \searrow G & \\ \mathcal{M}/T & \xrightarrow{\hat{T}} & [\text{Set}, \text{Set}] \end{array} \quad \cong$$

(more precisely, it is pseudoterminal in the 2-category of \mathcal{M} -gradings of T)

Example: \mathcal{M} = componentwise injective

\mathcal{M} -gradings of monads

For a class \mathcal{M} of natural transformations \succrightarrow , an \mathcal{M} -grading of a monad T consists of

- ▶ a monoidal category \mathcal{E}
- ▶ a graded monad $G : \mathcal{E} \rightarrow [\text{Set}, \text{Set}]$
- ▶ a monoidal natural transformation $m_e : Ge \succrightarrow T$, whose components are in \mathcal{M}

Under some conditions, the canonical \mathcal{M} -grading of the functor T gives a canonical \mathcal{M} -grading of the monad T

$$\begin{array}{ccc} \mathcal{E} & & \\ \downarrow F & \searrow G & \\ \mathcal{M}/T & \xrightarrow{\hat{T}} & [\text{Set}, \text{Set}] \end{array} \quad \cong$$

(more precisely, it is pseudoterminal in the 2-category of \mathcal{M} -gradings of T)

\mathcal{M} -gradings of monads

If

- ▶ \mathcal{M} forms a factorization system $(\mathcal{E}, \mathcal{M})$ on $[\mathbf{Set}, \mathbf{Set}]$

$$\begin{array}{ccc}
 F & \xrightarrow{f} & G \\
 e \searrow & & \nearrow m \\
 & S &
 \end{array}
 \quad (e \in \mathcal{E}, m \in \mathcal{M})$$

- ▶ for each $e : F \rightarrow F'$ in \mathcal{E} and G ,

$$(e \cdot G) : F \cdot G \rightarrow F' \cdot G \quad (G \cdot e) : G \cdot F \rightarrow G \cdot F'$$

are both in \mathcal{E}

then we get a canonical \mathcal{M} -grading of T , using

$$\begin{array}{ccc}
 \text{Id} & \xrightarrow{\eta} & T \\
 \searrow & & \nearrow \\
 & J &
 \end{array}
 \quad
 \begin{array}{ccc}
 S_1 \cdot S_2 & \xrightarrow{\quad} & T \cdot T \xrightarrow{\mu} T \\
 \searrow & & \nearrow \\
 & S_1 \boxtimes S_2 &
 \end{array}$$

$$JX = \{\eta x \mid x \in X\} \quad (S_1 \boxtimes S_2)X = \{\mu t \mid t \text{ is in the image of } S_1(S_2X) \hookrightarrow T(TX)\}$$

We can do something similar if we replace $[\mathbf{Set}, \mathbf{Set}]$ with some other monoidal category

Grading by sets of shapes

Every monad T on \mathbf{Set} has a set of shapes $T1$, and every $t \in TX$ has a shape $T!t \in T1$

- ▶ e.g. $\text{List}1 \cong \mathbb{N}$, and $T!t$ is the length of $t \in \text{List } X$

Whenever T is cartesian, subsets $\Sigma \subseteq T1$ form a canonical cartesian grading of T :

- ▶ subsets $\Sigma \subseteq T1$ are equivalently cartesian subfunctors $S \subseteq T$, via

$$\Sigma = S1 \quad SX = \{t \in TX \mid T!t \in \Sigma\}$$

(S is automatically cartesian)

families of subsets $SX \subseteq TX$
closed under $Tf : TX \rightarrow TY$
satisfying $SX = \{t \mid T!t \in S1\}$

- ▶ up to isomorphism, cartesian subfunctors form a factorization system $(\mathcal{E}, \mathcal{M})$ on $[\mathbf{Set}, \mathbf{Set}]_{\text{cart}}$

$$\begin{array}{ccc} F1 & \xrightarrow{\quad} & G1 \\ & \searrow & \nearrow \\ & S1 & \end{array}$$

Grading by sets of shapes

When $T = \text{List}$:

- ▶ Cartesian subfunctors $S \hookrightarrow \text{List}$ are equivalently subsets

$$\Sigma \subseteq \mathbf{N}$$

(allowable lengths of lists)

- ▶ which form an ordered monoid with

$$\Sigma \leq \Sigma' \Leftrightarrow \Sigma \subseteq \Sigma' \quad \mathbf{J} = \{0\} \quad \Sigma_1 \boxplus \Sigma_2 = \{n_1 + n_2 \mid n_1 \in \Sigma_1, n_2 \in \Sigma_2\}$$

- ▶ The canonical cartesian grading of List is

$$\text{List}'\Sigma X = \{xs \in \text{List}X \mid \text{length } xs \in \Sigma\}$$

Algebraic operations

An algebraic operation for T is a function

[Plotkin and Power '03]

$$T^n \rightarrow T$$

compatible with the multiplication of T

Given subfunctors S_1, \dots, S_n of T , we get

- ▶ a canonical subfunctor S'
- ▶ a flexibly graded algebraic operation

$$\hat{T}(S_1 \boxtimes -) \times \cdots \times \hat{T}(S_n \boxtimes -) \rightarrow \hat{T}(S' \boxtimes -)$$

(under some extra conditions about products)

Summary

Given a suitable class \mathcal{M} of natural transformations, every monad T has a **canonical \mathcal{M} -grading** $\hat{T} : \mathcal{M}/T \rightarrow [\mathbf{C}, \mathbf{C}]$

$$\begin{array}{ccc} \mathcal{E} & & \\ \downarrow F & \searrow G & \\ \text{Sub}(T) & \xrightarrow{\hat{T}} & [\mathbf{C}, \mathbf{C}] \end{array}$$

In particular, each monad T on Set has canonical gradings by

- ▶ subfunctors $S \hookrightarrow T$
- ▶ subsets $S1 \subseteq T1$ (assuming T is cartesian)